

On the Orientable Genus of Graphs Embedded in the Klein Bottle

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ABSTRACT

Let G be a graph embedded in the Klein bottle with "representativity" at least four. We give a formula for the orientable genus of G , which also implies a polynomially bounded algorithm. The formula is in terms of the number of times certain closed curves on the Klein bottle intersect the graph. In particular, it shows that a cut-and-paste technique for re-embedding graphs is the best possible.

1. INTRODUCTION

Let ψ be an embedding of a graph G in a surface Σ . We say that ψ is an embedding with *representativity* ρ if ρ is maximum such that every nonnull-homotopic closed curve intersects the image of G at least ρ times. The following is a result of [2].

1.1. Let G be a graph embedded in the projective plane with representativity $\rho \geq 3$. Then the orientable genus of G is $\lfloor \rho/2 \rfloor$.

In this paper we prove an analogue of 1.1 for the Klein bottle, which also yields a polynomially bounded algorithm. Testing genus is NP-complete for general graphs [5], and therefore it would be interesting to characterize all

lower ideals for which testing genus is polynomial-time solvable. Our result can be viewed as a step towards this problem.

If X is a topological space, a *closed curve* (in X) is a continuous mapping $f: S \rightarrow X$, where S is the unit circle in the complex plane \mathbb{C} . A closed curve is called *simple* if it is 1 - 1. The Klein bottle, denoted by N_2 , is the space obtained by identifying sides a, b and c, d of a square with sides a, b, c, d as shown in Figure 1. Then the identified sides a, b form a closed curve α_1 , and the identified sides c, d form a closed curve α_2 ; we may assume that $\alpha_1(1) = \alpha_2(1)$. Given two closed curves α, β with $\alpha(1) = \beta(1)$, we define a closed curve $\alpha\beta$ by

$$\alpha\beta(e^{it}) = \begin{cases} \alpha(e^{2it}), & 0 \leq t < \pi; \\ \beta(e^{2it}), & \pi \leq t < 2\pi. \end{cases}$$

Let $x_0 = \alpha_1(1) = \alpha_2(1)$. We recall that the fundamental group $\pi(N_2, x_0)$ has a presentation in terms of generators $\{\alpha_1, \alpha_2\}$ and one relation $\alpha_1^2\alpha_2^2$. For $\gamma \in \pi(N_2, x_0)$ we say that a closed curve f in N_2 is a γ -curve, if f is (freely) homotopic to γ or γ^{-1} . (Warning: sometimes a γ_1 -curve can be homotopic to a γ_2 -curve for $\gamma_1, \gamma_2 \in \pi(N_2, x_0)$ such that $\gamma_1 \neq \gamma_2$ and $\gamma_1 \neq \gamma_2^{-1}$). Thus an α_1^2 -curve (same as α_2^2 -curve) separates the Klein bottle into two Möbius bands.

By a *surface* we mean a compact connected 2-manifold with (possibly null) boundary. Let (f_1, \dots, f_n) be a family of closed curves in a surface Σ . We put $M(f_1, \dots, f_n) = \{(i, z), (i', z') : (i, z) \neq (i', z'), i, i' \in \{1, \dots, n\}, z, z' \in S \text{ and } f_i(z) = f_{i'}(z')\}$. The family (f_1, \dots, f_n) is said to be *finitary* if $M(f_1, \dots, f_n)$ is finite. For a finitary family (f_1, \dots, f_n) of closed curves we define $\text{cr}(f_1, \dots, f_n)$, the *crossing number*, as the cardinality of the set

$$\{(i, z), (i', z')\} \in M(f_1, \dots, f_n); f_i, f_{i'} \text{ cross at } (z, z'),$$

where closed curves f, h are said to cross at (z, z') if $f(z) = h(z')$ and there exists a homeomorphism of a neighborhood of $f(z) = h(z')$ onto the unit disk in \mathbb{C} such that the image of f follows the real axis and the image of h follows the imaginary axis. The family (f_1, \dots, f_n) is said to be *cross-free* if $\text{cr}(f_1, \dots, f_n) = 0$.

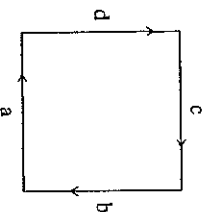


FIGURE 1

A *1-blockage* is a finitary 1-element family (f) consisting of an α_1 - α_2 -curve f in N_2 . A *2-blockage* is a finitary (ordered) pair (f_1, f_2) of two closed curves in N_2 such that f_i is an α_i -curve ($i = 1, 2$). A *blockage* is a 1-blockage or a 2-blockage.

Let G be a graph embedded in a surface Σ . The G -*degree* (or simply *degree* when no confusion is likely) of a closed curve f in Σ , denoted by $\text{deg } f$, is the cardinality of the set $\{z \in S; f(z) \in G\}$. Now let $\Sigma = N_2$. The *order* of a 1-blockage (f) is $\lceil \text{deg } f/2 \rceil$, the order of a 2-blockage (f_1, f_2) is $\lfloor \text{deg } f_1/2 \rfloor + \lfloor \text{deg } f_2/2 \rfloor$. We can state our main result now.

1.2. Let G be a graph embedded in N_2 in such a way that every α_i -curve has degree at least four ($i = 1, 2$), and let $g \geq 4$ be an integer. Then the following conditions are equivalent.

- (i) G has orientable genus $\geq g$.
- (ii) Every cross-free blockage has order $\geq g$.
- (iii) Every blockage has order $\geq g$.

Since (iii) above is easy to test in polynomial time (hint: consider a lifting of G in the universal covering surface), we have the following algorithmic result.

1.3. There exists a polynomial algorithm to solve the following problem.
 Instance: A graph G , embedded in N_2 as in (1.2), and an integer $g \geq 0$.
 Question: Is the orientable genus of G at least g ?

For if $g = 0, 1, 2$, or 3 , we can use the algorithm of [3]. If $g \geq 4$ we apply (1.2).

The assumption about the degree of every α_i -curve in (1.2) is necessary, because in the embedding of a graph G shown in Figure 2 every blockage has order ≥ 4 and yet Figure 3 shows the embedding of G in S_3 , the orientable surface without boundary of genus 3. In fact, it follows from our proof of 1.2 that even if the embedding of G into N_2 does not satisfy the assumption about the degree of every α_i -curve, then the orientable genus of G can only be either g or $g - 1$. We have found a polynomial-time algorithm to decide whether it is g or $g - 1$ (notice that the existence of such an algorithm does not seem to follow from [3]), but the algorithm is quite complicated and does not seem to contain any essentially new ideas, and so is not described here.

We conjecture the following generalization of 1.2. Let N_k be the nonorientable surface with cross-cap number k , that is, the surface obtained by identifying k pairs of consecutive sides of a $2k$ -gon. A blockage (in N_k) is a finitary n -tuple $B = (f_1, \dots, f_n)$ of closed curves such that there exist closed curves f'_1, f'_2, \dots, f'_n with the property that f_i is homotopic to f'_i

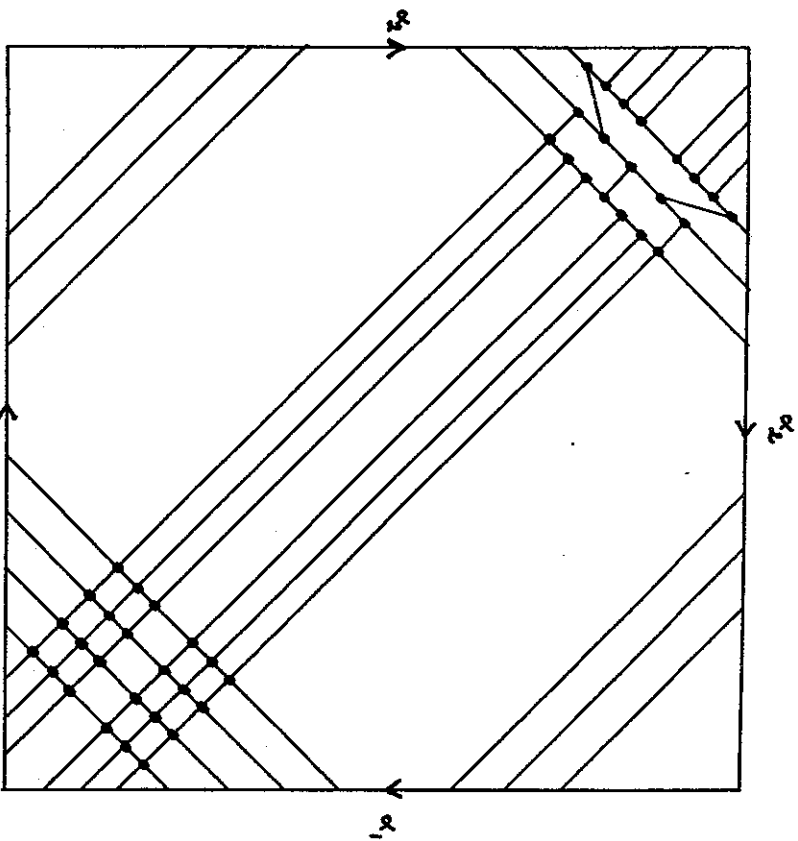


FIGURE 2

($i = 1, \dots, n$), $M(f_1^i, \dots, f_n^i) = \emptyset$ and the surface Σ obtained by cutting open along all of f_1, f_2, \dots, f_n is orientable. We define $\text{res}(B)$, the *residuum* of B , to be the orientable genus of the surface Σ . We define the order of f_i ($i = 1, \dots, n$) by

$$\text{ord}(f_i) = \begin{cases} \frac{\deg f_i}{2}, & \text{if } f_i \text{ is orientation reversing,} \\ \frac{\deg f_i}{2}, & \text{if } f_i \text{ is orientation preserving,} \end{cases}$$

where a curve g is said to be *orientation reversing* (*orientation preserving*, respectively) if passing once around g changes (does not change, respectively) the orientation. Finally, we define the order of B as

$$\text{res}(B) + \sum_{i=1}^n \text{ord}(f_i).$$

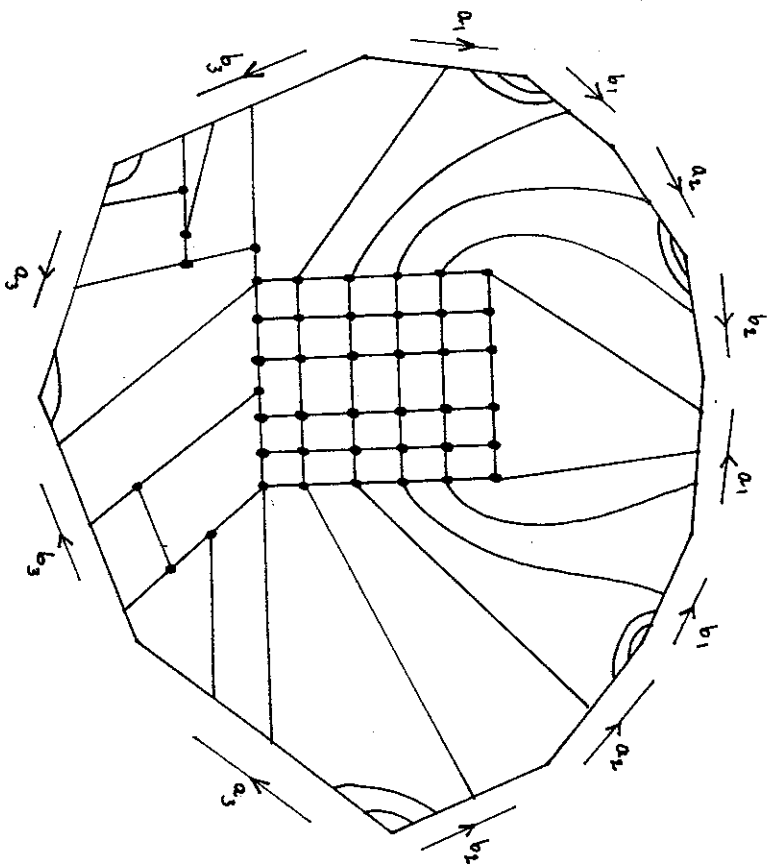


FIGURE 3

1.4. Conjecture. Let $k \geq 2$ be a fixed integer and let G be a graph embedded in N_g with sufficient representativity, and let $g \geq 2k$ be an integer. Then the following conditions are equivalent:

- (i) The orientable genus of G is at least g .
- (ii) Every cross-free blockage has order at least g .
- (iii) Every blockage has order at least g .

We end this section with some more terminology. A *graph* is a pair $G = (U, V)$, where U is a Hausdorff topological space and $V \subseteq U$ is a finite set such that

- (i) $U - V$ is the disjoint union of finitely many open subsets, called *edges*,
- (ii) for each edge e , its boundary is a subset of V consisting of one or two points. If it consists of two points, then the pair (\bar{e}, e) is homeomorphic to the pair $([0, 1], (0, 1))$; if the boundary consists of one point, then the pair (\bar{e}, e) is homeomorphic to the pair $(S, S - \{1\})$,

(iii) a subset $A \subseteq U$ is closed (open) if and only if $A \cap \bar{e}$ is closed (open) for all edges e .

[Here \bar{e} denotes the closure of the set e .] We write $U(G) = U, V(G) = V$, and we denote the set of edges by $E(G)$. The rest of our graph terminology is standard, we remark that paths and circuits have no "repeated" vertices and no "repeated" edges. We say that a mapping ψ is an *embedding* of a graph G into a topological space X if $\psi: U(G) \rightarrow X$ is a homeomorphism onto its image. We say that a graph G is *embedded* in a topological space X if the identity is an embedding of G into X . Obviously, if ψ is an embedding of G into X , then the image graph $(\psi(U(G)), \psi(V(G)))$ is isomorphic to G and is embedded in X . Therefore we shall usually assume that the graph in question is actually embedded in X .

We say that closed curves f, g in X are *homotopic* if there exists a continuous function $F: S \times [0, 1] \rightarrow X$ such that $F(z, 0) = f(z)$ and $F(z, 1) = g(z)$ for every $z \in S$. A closed curve f is *null-homotopic* if it is homotopic to some constant curve, and is *nonnull-homotopic* otherwise.

If $f: X \rightarrow Y$ is a mapping then $f(X)$ denotes the set $\{f(x): x \in X\}$. Let G be a graph embedded in N_2 and let f be a simple γ -curve for some $\gamma \in \pi(N_2, x_0)$ such that $f(S)$ is a circuit in G . We say that $f(S)$ is a γ -*circuit* in G .

2. DISJOINT PATHS IN THE KLEIN BOTTLE

In this section we prove three results that we shall need for the proof of 1.2 in Section 3. We start with a well-known and easy lemma.

2.1. Let f_1, f_2 be two nonnull-homotopic closed curves in N_1 , the projective plane. Then $f_1(S) \cap f_2(S) \neq \emptyset$.

A *pseudocylinder* is a surface obtained from a sphere by deleting interiors of two regions R_1, R_2 homeomorphic to closed discs with $R_1 \cap R_2$ finite. We refer to the boundaries of R_1 and R_2 as the boundary components.

2.2. Let $k \geq 1$ be an integer and let G be a graph embedded in a pseudocylinder Σ . Then either there exists a simple nonnull-homotopic closed curve f in Σ of degree $< k$ separating the two boundary components, or else there are k disjoint paths in G , each joining the two boundary components of Σ .

Proof: This follows easily from Menger's theorem. ■

The following is the first result of this section.

2.3. Let $k \geq 0$ be an integer and let G be a graph embedded in N_2 in such a way that every cross-free 1-blockage has order $\geq k$. Then there

exist disjoint circuits C_1, \dots, C_k in G such that C_i separates C_{i-1} and C_{i+1} ($i = 2, \dots, k - 1$) and one of the following conditions is satisfied:

- (a) C_1 is an α_1 -circuit and C_k is an α_1^2 -circuit ($i = 2, \dots, k$);
- (b) C_1 is an α_2 -circuit and C_k is an α_1^2 -circuit ($i = 2, \dots, k$);
- (c) C_i is an α_i^2 -circuit ($i = 1, \dots, k$).

Proof. Let (h) be a cross-free 1-blockage with $n = |\{z \in S: h(z) \in U(G)\}|$ minimum. Then $n \geq 2k - 1$. Let v_1, \dots, v_n (in this order) be the interior sections of h with $U(G)$. We cut open along h and thus obtain an embedding of a graph G' in a pseudocylinder Σ with, say, vertices x_1, \dots, x_n (in this order) on one boundary component, and say vertices y_n, \dots, y_1 (in this order) on the other boundary component, so that G' is obtained from G' by identifying x_i and y_i into v_i ($i = 1, \dots, n$). There is no simple closed curve h' in Σ separating the two boundary components of G' -degree $< n$, because after pasting the two boundary components together, (h') becomes a 1-blockage in N_2 . Hence by 2.2 there are n disjoint paths in G' , each joining x_i and y_{i_0} ($i = 1, \dots, n$). It follows that $j(j(i)) = i$, and therefore, after pasting the two boundary components together, the paths P_1, \dots, P_n become the desired k disjoint circuits. ■

The proof of the following lemma is left to the reader.

2.4. Let G be a graph embedded in N_1 , the projective plane, with representativity ≥ 1 . Then there exists a nonnull-homotopic simple closed curve h in N_1 such that $h(S)$ is a circuit in G .

Let X be a topological space. A curve (in X) is a continuous mapping $f: I \rightarrow X$, where I is the closed unit interval $[0, 1]$.

Let f be a null-homotopic simple closed curve in N_1 . Then f bounds a disk with interior U . Let $\Sigma = N_1 - U$. Then Σ is a *Möbius band* and f is called a *boundary curve* of Σ . A curve $h: I \rightarrow \Sigma$ is called a *chord* (in Σ) if $h(0)$ and $h(1)$ belong to the boundary of Σ . A chord $h: I \rightarrow \Sigma$ is called *essential* if the closed curve g is nonnull-homotopic in N_1 , where

$$g(e^{it}) = \begin{cases} h((t - t_0)/(t_1 - t_0)), & \text{if } t_0 \leq t < t_1, \\ f(e^{it}), & \text{if } t_1 \leq t < t_0 + 2\pi, \end{cases}$$

and t_0, t_1 are such that $f(e^{it_0}) = h(0), f(e^{it_1}) = h(1), |t_0 - t_1| \leq 2\pi$ and $t_0 < t_1$. The following lemma follows easily from 2.1.

2.5. Let Σ be the Möbius band and let h_1, h_2 be two essential chords with $h_1(I) \cap h_2(I) = \emptyset$. Then $h_1(0), h_2(0), h_1(1)$, and $h_2(1)$ occur on the boundary of Σ in this (or reverse) order.

2.6. Let G be a graph embedded in the Möbius band Σ with distinct vertices u, v on the boundary of Σ . Then either there exists a closed curve of

degree ≤ 1 , homotopic to the boundary curve or else there exists an essential chord h in Σ such that $h(0) = u$, $h(1) = v$, and $h(I)$ is a path in G .

Proof. Let G be as stated and assume that every closed curve homotopic to the boundary curve has degree ≥ 2 . Let G' be the graph obtained from G by identifying the vertices u and v ; the embedding of G induces an embedding of G' in N_1 . By 2.4 there exists a nonnull-homotopic simple closed curve f in N_1 such that $f(S)$ is a circuit in G' , let C be the circuit or path generated by $f(S)$ in G . It follows from the assumption that there exist two disjoint paths in G from $\{u, v\}$ to C , combining these paths with a suitable segment of C yields the desired chord. ■

The following is the second result.

2.7. Let G be a graph embedded in N_2 in such a way that every simple closed α_1^2 -curve has degree at least 2. Then there exists an $\alpha_1\alpha_2$ -circuit in G .

Proof. Suppose, if possible, that the theorem fails and let G be a counterexample with $|V(G)|$ minimum. Then there exists a simple α_1^2 -curve h with degree 2, let u, v be the intersections of h with $U(G)$. We may assume that $u, v \in V(G)$. When we cut open along h we obtain two Möbius bands Σ_1, Σ_2 . By 2.6 there exists an essential chord h_i in Σ_i with $h_i(0) = u$, $h_i(1) = v$, and such that $h_i(I)$ is a path in G ($i = 1, 2$). It is easily seen that $h_1(I) \cup h_2(I)$ is as desired. ■

For an orientable surface, the last result of this section would follow from a theorem of Schrijver [4]. However, his theorem is still open for nonorientable surfaces. We say that curves $f, g: I \rightarrow N_2$ are *equivalent* if there exists a continuous mapping $F: I \times I \rightarrow N_2$ such that $F(t, 0) = f(t)$, $F(t, 1) = g(t)$, $F(0, t) = f(0)$ and $F(1, t) = f(1)$ for all $t \in I$.

2.8. Let f_i be a simple α_i -curve in N_2 ($i = 1, 2$), assume that f_1, f_2 cross each other, and that (f_1, f_2) is finitary. Then there exist t_1, t_2, t'_1, t'_2 with $|t_1 - t'_1| < 2\pi$, $|t_2 - t'_2| < 2\pi$, $t_1 < t'_1, t_2 < t'_2$ such that f_1, f_2 cross at (e^{it_1}, e^{it_2}) and at $(e^{it'_1}, e^{it'_2})$ and such that if we define, for $t \in [0, 1]$

$$\begin{aligned} g_1(t) &= f_1(e^{it(1-0)+t(1)}), \\ g_2(t) &= f_1(e^{it(t-0)+t(1-2\pi)}), \\ h_1(t) &= f_2(e^{it(1-0)+t(1/2)}), \\ h_2(t) &= f_2(e^{it(2(1-0)+t(1/2-2\pi))}), \end{aligned}$$

then there exist $i, j \in \{1, 2\}$ such that g_i is equivalent to h_j .

Proof. Let $u_1, \dots, u_n \in N_2$ be the points where f_1 and f_2 cross. We cut open along f_1 and obtain a Möbius band Σ , with say, vertices $x_1, \dots, x_n, y_1, \dots, y_n$ (in this order) on the boundary, where x_i, y_i are obtained by split-

ting u_i . The segments of f_2 become disjoint chords in Σ . By 2.5 at least one of these chords is not essential (otherwise each chord would join x_i and y_i , which is impossible, since f_2 is a closed curve). The theorem now follows from the definition of an essential chord. ■

3. MAIN RESULT

We shall prove 1.2 by proving (i) \Rightarrow (ii) \Rightarrow (iii) \Rightarrow (i). We denote by S_g the orientable surface without boundary of genus g .

3.1. Let G be a graph embedded in N_2 . If there exists a cross-free blockage of order $\leq g$, then the orientable genus of G is at most g .

Proof. We prove the result by an elementary technique called "surgery." Let B be a cross-free blockage of order $\leq g$. We first assume that B is a 1-blockage. Let $B = (f)$, and let v_1, \dots, v_k ($k \leq 2g$) be the intersections (in this order) of f with $U(G)$. We cut open along f and split each v_i into two vertices x_i and y_i , to obtain another graph G' embedded into a pseudocylinder Σ with the vertices x_1, \dots, x_k (in this order) on one boundary component, and the vertices y_1, \dots, y_k (in this order) on the other boundary component. It is easily seen that we can add $\leq g$ handles to Σ , to obtain an orientable surface S_g of genus g , and identify vertices x_i and y_i ($i = 1, \dots, k$) along the added handles. Thus we can produce an embedding of G into S_g .

Now if B is a 2-blockage, say $B = (f_1, f_2)$, we proceed similarly. Let v_1, \dots, v_k be the intersections of f_1 with $U(G)$, and let v'_1, \dots, v'_k be the intersections of f_2 with $U(G)$. We first cut open along f_1 and split each v_i into x_i and y_i , and then cut open along f_2 and split each v'_i into x'_i and y'_i . We obtain a graph G' embedded into a pseudocylinder with vertices $x_1, \dots, x_k, y_1, \dots, y_k$ (in this order) on one boundary component, and with vertices $x'_1, \dots, x'_k, y'_1, \dots, y'_k$ (in this order) on the other boundary component. Again, by adding at most $\lfloor k/2 \rfloor + \lfloor k'/2 \rfloor \leq g$ handles, we can identify x_i and y_i ($i = 1, \dots, k$), and x'_i and y'_i ($i = 1, \dots, k'$) to obtain an embedding of G into S_g . ■

3.2. Let G be a graph embedded in N_2 as in 1.2 and let B be a blockage of order $g \geq 0$. Then there exists a cross-free blockage of order $\leq g$.

Proof. Let us choose a blockage B of order $\leq g$ with $\text{cr}(B)$ minimum, and subject to that with $|M(B)|$ minimum. We suppose that $\text{cr}(B) > 0$ and seek a contradiction. Let $B = (f_1, \dots, f_n)$ (so that $n = 1$ or 2) and let H be the pair $(f_1(S), \{f_i(z) : \{1, z\}, (1, z')\} \in M(f_1)$ for some $z' \in S\})$. Then H is a graph unless f_1 is simple.

(1) If f_1 is not simple then every α_1 -curve has H -degree ≥ 1 , and if $n = 1$ then every α_2 -curve has H -degree ≥ 1 .

For this follows from 2.1.

(2) If f_1 is not simple then there exists a simple α_1^2 -curve h in N_2 with H -degree at most 1.

For otherwise there exists by 2.7 a simple $\alpha_1\alpha_2$ -curve f in N_2 with $f(S) \subseteq U(H)$. Then (f) is a cross-free blockage of order $\leq g$ (this is obvious if $n = 1$, and if $n = 2$ then it follows from the fact that f_2 has G -degree ≥ 4), a contradiction.

(3) $n = 2$ and f_1, f_2 are simple.

For suppose that either $n = 1$ (in which case f_1 is not simple because $\text{cr}(B) > 0$), or that f_1 is not simple. The curve h from (2) separates N_2 into two Möbius bands Σ_1, Σ_2 . We may assume that the notation is chosen so that Σ_1 contains an α_1 -curve. By (1) and 2.4 there exists a simple α_1 -curve f'_1 with $f'_1(S) \subseteq f_1(S) \cap \Sigma_1$. If $n = 1$ then similarly there exists a simple α_2 -curve f'_2 with $f'_2(S) \subseteq f_1(S) \cap \Sigma_2$, but then (f'_1, f'_2) is a cross-free blockage of order $\leq g$, a contradiction. If $n = 2$ then $B' = (f'_1, f_2)$ is a blockage of order $\leq g$ with $\text{cr}(B') \leq \text{cr}(B)$ and $|M(B')| < |M(B)|$, a contradiction. Hence $n = 2$ and f_1 is simple. We deduce from the symmetry that f_2 is also simple. This proves (3).

We deduce that f_1 and f_2 cross each other and so we may apply 2.8. We may assume without loss of generality that (using the notation of 2.8) $t_1 = t_2, t'_1 = t'_2$, and that g_1 and h_1 are equivalent. We may also assume that

$$|\{t \in (0, 1) : g_1(t) \in U(G)\}| \leq |\{t \in (0, 1) : h_1(t) \in U(G)\}|.$$

Let

$$f_2^i(e^{it}) = \begin{cases} f_1(e^{it}), & \text{if } t_1 \leq t < t'_1; \\ f_2(e^{it}), & \text{if } t'_1 \leq t < t_1 + 2\pi. \end{cases}$$

It follows that f_2^i is homotopic to f_2 and that its G -degree is at most the G -degree of f_2 . We can shift f_2^i slightly into a closed curve f_2'' of the same G -degree and homotopic to f_2 such that (f_1, f_2'') is finitary and $\text{cr}(f_1, f_2'') < \text{cr}(f_1, f_2)$. This is a contradiction, because (f_1, f_2'') is a blockage of order $\leq g$. ■

We require the following result of Battle, Harary, Kodama, and Youngs [1].

3.3. Let the graph G be a disjoint union of graphs G_1 and G_2 . Then the orientable genus of G is $g_1 + g_2$, where g_i is the orientable genus of G_i ($i = 1, 2$).

3.4. Let G_1, G_2 be two disjoint connected graphs of genera $g_1, g_2 \geq 1$, respectively, embedded in $S_{g_1+g_2}$. Then there exists a simple closed curve h in $S_{g_1+g_2}$ separating the surface into two surfaces Σ_1, Σ_2 with common boundary in such a way that Σ_1 contains G_1 and Σ_2 contains G_2 (and therefore the orientable genus of Σ_i is g_i).

Proof. There exist simple closed curves h_1, \dots, h_n in $S_{g_1+g_2}$ such that $\cup_{i=1}^n h_i(S)$ separates the two graphs. Let us choose such a system with n minimum. It follows that none of the curves is null-homotopic and by 3.3 each separates the surface. Hence $n = 1$ and the partition is as desired, again by 3.3. ■

3.5. Let G be a graph embedded in N_2 as in 1.2 and assume that every blockage has order $\geq g$, where $g \geq 4$ is some integer. Then G has orientable genus $\geq g$.

Proof. Suppose for a contradiction that G is a graph that satisfies the hypothesis but not the conclusion of the theorem with $|V(G)|$ minimum. We shall assume that (a) of 2.3 occurs, for the case (b) is symmetric and (c) is handled similarly and is in fact easier. So let C_1, \dots, C_g be as in 2.3(a). Let $\rho_1 \geq 4$ and $\rho_2 \geq 4$ be such that $\lfloor \rho_1/2 \rfloor + \lfloor \rho_2/2 \rfloor = g$ and such that every α_i -curve in N_2 has degree at least ρ_i . Let us choose a simple α_1^2 -curve f with the following properties:

- (i) $\text{Im } f$ separates $C_{\lfloor \rho_1/2 \rfloor}$ and $C_{\lfloor \rho_2/2 \rfloor+1}$.
- (ii) $\text{Im } f \cap V(G) = \emptyset$.
- (iii) $|\text{Im } f \cap U(G)|$ is minimum subject to (i) and (ii).

Let E be the set of edges intersected by $\text{Im } f$, and let G_0 be the graph obtained from G by deleting all edges from E . The curve f separates N_2 into two Möbius bands Σ_1, Σ_2 such that Σ_1 contains C_1 , and it separates G_0 into two disjoint subgraphs G_1 and G_2 . Let G_1 be the one that contains C_1 . Let $\Sigma'_i = \Sigma_i - U_i$, where Σ'_i is homeomorphic to N_1 and U_i is an open disk in Σ'_i . Then the embedding of G into N_2 induces an embedding of G_i into Σ'_i ($i = 1, 2$).

- (1) The orientable genus of G_1 is at least $\lfloor \rho_1/2 \rfloor - 1$.

For we consider the embedding of G_1 in Σ'_1 . Let h be a nonnull-homotopic closed curve in Σ'_1 , then h can be regarded as an α_1 -curve in N_2 ; and if h intersects an edge from E , or if it intersects G_2 , then it meets C_1 at least once, and each of $C_2, \dots, C_{\lfloor \rho_1/2 \rfloor}$ at least twice. Since h has G -degree $\geq \rho_1$, it follows that it has G_1 -degree $\geq 2\lfloor \rho_1/2 \rfloor - 1$. Hence G_1 is embedded in Σ'_1 with representativity $\geq 2\lfloor \rho_1/2 \rfloor - 1$ and thus its orientable genus is $\geq \lfloor \rho_1/2 \rfloor - 1$ by 1.1.

- (2) The orientable genus of G_2 is at least $\lfloor \rho_2/2 \rfloor$.

For this follows similarly as (1).

(3) The graphs G_1 and G_2 are connected.

For let $i \in \{1, 2\}$. We consider G_i as being embedded in Σ_i . By the argument of (1) and 2.4, G_i contains a nonnull-homotopic circuit; let G_i' be the component of G_i containing this circuit. By 2.1 every other component B of G_i contains no nonnull-homotopic circuit and is therefore contained in a disk. It follows from the choice of f that B is a component of G , which is impossible by the minimality of $|V(G)|$.

Let G' be the graph obtained from G by contracting G_2 to a vertex v_0 .

(4) The orientable genus of G' is at least $\lfloor \rho_1/2 \rfloor$.

For the embedding of G_1 in Σ_1' has the property that all endpoints of edges of E that are in $V(G_1)$ lie on the boundary of the face containing U_1 . Therefore we can construct an embedding ψ' of G' into Σ_1' by choosing $\psi'(v_0) \in U_1$ and $\psi'(e)$ for $e \in E(G') - E(G_1)$ appropriately. We claim that the representativity of this embedding is at least $2\lfloor \rho_1/2 \rfloor$. For suppose, if possible, that h is a nonnull-homotopic simple closed curve in Σ_1' of G' -degree $< 2\lfloor \rho_1/2 \rfloor \leq \rho_1$. Then h intersects neither v_0 , nor any edge in E , for otherwise it would have to intersect each C_i ($i = 2, \dots, \lfloor \rho_1/2 \rfloor$) at least twice and C_1 at least once, which would be $\geq 2\lfloor \rho_1/2 \rfloor$ intersections altogether. But now it follows that h can be regarded as an α_1 -curve in N_2 of G -degree $< \rho_1$, a contradiction. Hence no such h exists and (4) follows from 1.1.

Now suppose for a contradiction that there exists an embedding ψ of G into S_{g-1} . By (1), (2), (3), and 3.4 there exists a simple closed curve h in S_{g-1} separating S_{g-1} into two surfaces Σ_1' and Σ_2' such that Σ_1' has genus $\lfloor \rho_1/2 \rfloor - 1$ and contains the image of G_1 . Hence G_1 has an embedding ψ_1 into Σ_1' with the property that all end points of edges of E that belong to $V(G_1)$ lie on the boundary of the face incident with $h(S)$. Similarly, as in the proof of (4) we deduce that G' has an embedding into $S_{\lfloor \rho_1/2 \rfloor - 1}$, contrary to (4). ■

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