## Szemerédi's regularity lemma.

Let  $\varepsilon > 0$ . Let  $A, B \subseteq V(G)$  be disjoint. We say that (A, B) is  $\varepsilon$ -**regular** in *G* if for all  $X \subseteq A$  and  $Y \subseteq B$  with  $|X| \ge \varepsilon |A|$  and  $|Y| \ge \varepsilon |B|$  we have

$$|d(A,B) - d(X,Y)| \le \varepsilon$$

where

$$d(C,D) = \frac{|\langle C,D\rangle|}{|C|\cdot|D|}$$

and  $\langle C, D \rangle = \{e: e \text{ has one end in } C, \text{ the other in } D\}.$ 

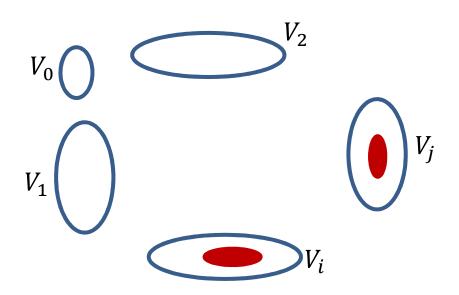
A partition  $(V_0, V_1, ..., V_k)$  of V(G) is  $\varepsilon$ -regular if

(i)  $|V_0| \le \varepsilon |V(G)|$ 

(ii)  $|V_1| = |V_2| = \dots = |V_k|$ 

(iii) all but at most  $\varepsilon k^2$  pairs  $(V_i, V_j)$  are  $\varepsilon$ -regular  $(1 \le i, j \le k)$ .

**Theorem** (Szemerédi's regularity lemma).  $\forall \varepsilon > 0 \forall$  integer  $m \exists$  integer  $M \forall$  graph G on  $n \ge m$  vertices  $\exists k$  with  $m \le k \le M$  and an  $\varepsilon$ -regular partition  $(V_0, V_1, V_2, ..., V_k)$  of V(G).



#### The Erdős-Stone theorem

# **Turán's theorem.** If G has no $K_r$ subgraph, then

$$|E(G)| \le |E(T_{r-1}(n))|,$$

with equality if and only if  $G \cong T_{r-1}(n)$ .

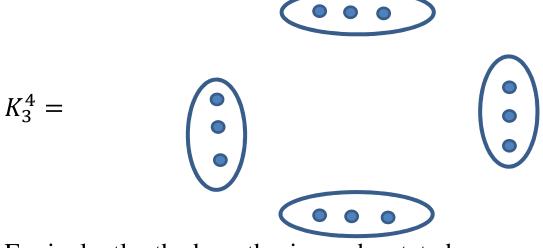
Recall that  $T_{r-1}(n) = \text{complete } (r-1)\text{-partite graph on } n$  vertices with color classes as close to each other in size as possible. Let

$$t_{r-1}(n) \coloneqq \left| E\left(T_{r-1}(n)\right) \right| \approx \frac{r-2}{r-1} \cdot \frac{n^2}{2}$$

**Theorem.** (Erdős-Stone)  $\forall r, s \forall \varepsilon > 0 \exists n_0 \forall \text{ graph } G \text{ on } \ge n_0$  vertices if

$$|E(G)| \ge t_{r-1}(n) + \varepsilon n^2,$$

then G has a  $K_s^r \coloneqq K_{s,s,\dots,s}$ -subgraph.



Equivalently, the hypothesis can be stated as

$$|E(G)| \ge \left(\frac{r-2}{r-1} + \varepsilon\right) \frac{n^2}{2}$$

# **Definition.**

 $ex(n, H) \coloneqq max\{|E(G)|: |V(G)| = n, G \text{ has no } H \text{ subgraph}\}$ Example.  $ex(n, K_r) = t_{r-1}(n)$ .
Corollary  $\lim_{k \to \infty} \frac{ex(n, H)}{k} - \frac{\chi(H) - 2}{k}$  for every H with  $\geq 1$  edge

**Corollary.**  $\lim_{n\to\infty} \frac{\exp(n,H)}{\binom{n}{2}} = \frac{\chi(H)-2}{\chi(H)-1}$  for every *H* with  $\ge 1$  edge **Proof.** Weekly exercise.

**Definition.** The upper density of an (infinite) graph *G* is

$$\limsup \left\{ \frac{|E(H)|}{\binom{|V(H)|}{2}} : H \subseteq G, H \text{ finite} \right\}$$

Corollary. The upper density of an infinite graph is

0, 
$$\frac{1}{2}$$
,  $\frac{2}{3}$ ,  $\frac{3}{4}$ , ..., or 1

**Proof.** Weekly exercise.

**Proof of Erdős-Stone.** Let *r*, *s* and  $\gamma > 0$  be given. Want to show  $\exists n_0$  such that if  $|V(G)| \ge n_0$ , then

$$|E(G)| \ge \left(\frac{1}{2}\frac{r-1}{r-2} + \gamma\right)n^2 \Rightarrow K_s^r \subseteq G$$

Let  $\varepsilon = \varepsilon(r, s, \gamma)$  and  $m = m(\gamma)$ , TBD.

Let *M* be as in Szemerédi's regularity lemma. Let *G* have at least *m* vertices. By Szemerédi's regularity lemma *G* has an  $\varepsilon$ -regular partition  $(V_0, V_1, \dots, V_k)$  with  $m \le k \le M$ .

Note 
$$(1 - \varepsilon) \frac{n}{k} \le |V_i| \le \frac{n}{k}$$
 for  $i = 1, 2, ..., k$ .

Let *R* be the "regularity graph" defined by  $V(R) \coloneqq \{1, 2, ..., k\}$  and  $i \sim j$  in *R* if  $(V_i, V_j)$  is  $\varepsilon$ -regular and the density  $d(V_i, V_j)$  is  $\geq \gamma$ .

**Claim 1.**  $\varepsilon$ , *m* can be chosen so that *R* has a  $K_r$  subgraph.

**Pf.** If not, then by Turán's theorem  $|E(R)| \le t_{r-1}(k) \le \frac{1}{2} \frac{r-2}{r-1} k^2$ . We have that |E(G)| is at most the sum of:

- # of edges incident with  $V_0$ , which is  $\leq \epsilon n^2$
- # edges in some  $V_i$ , which is  $\leq k \frac{1}{2} \left(\frac{n}{k}\right)^2 \leq \frac{1}{2} \frac{n^2}{m}$
- # of  $V_i V_j$  edges for  $i, j \in E(R)$ , which is  $\leq |E(R)| \left(\frac{n}{k}\right)^2 \leq \frac{1}{2} \frac{r-2}{r-1} n^2$
- # of  $V_i V_j$  edges for  $(V_i, V_j)$  not  $\varepsilon$ -regular, which is  $\leq \varepsilon k^2 \left(\frac{n}{k}\right)^2$
- # of  $V_i V_j$  edges when  $d(V_i, V_j) < \gamma$ , which is  $\leq {\binom{k}{2}} \gamma \left(\frac{n}{k}\right)^2$

$$\begin{split} |E(G)| &\leq \varepsilon n^2 + k \frac{1}{2} \left(\frac{n}{k}\right)^2 + |E(R)| \left(\frac{n}{k}\right)^2 + \varepsilon k^2 \left(\frac{n}{k}\right)^2 + \left(\frac{k}{2}\right) \gamma \left(\frac{n}{k}\right)^2 \\ &\leq \frac{1}{2} \frac{r-2}{r-1} n^2 + \varepsilon n^2 + \frac{1}{2} \frac{n^2}{k} + \varepsilon n^2 + \frac{\gamma}{2} n^2 \leq \\ &\leq \frac{1}{2} \frac{r-2}{r-1} n^2 + \left(2\varepsilon + \frac{1}{2m} + \frac{\gamma}{2}\right) n^2 \end{split}$$

So on choosing  $\varepsilon$ , *m* such that  $4\varepsilon + \frac{1}{m} < \gamma$  we get a contradiction. This proves Claim 1.

Next we show that if  $\varepsilon$  is sufficiently small, then  $K_s^r \subseteq G$ . We may assume that  $\{1, 2, ..., r\}$  is a clique in *R*.

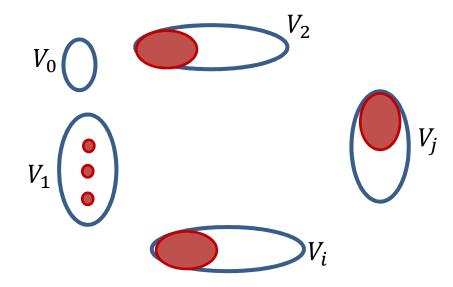
For  $1 \le i < j \le r$  we have  $d(V_i, V_j) \ge \gamma$  and so  $d(X, Y) \ge \gamma - \varepsilon$  for every  $X \subseteq V_i$  and  $Y \subseteq V_j$  with  $|X| \ge \varepsilon |V_i|$  and  $|Y| \ge \varepsilon |V_j|$ .

We will show:

(\*) Let  $V'_i \subseteq V_i$  satisfy  $|V'_i| \ge \varepsilon |V_i|$  and  $|V'_1| \ge (r-1)\varepsilon |V_1| + s$ . Then there exist distinct vertices  $v_1, v_2, \dots, v_s \in V'_1$  and sets  $V''_i \subseteq V'_i$  such that

- { $v_1$ , ...,  $v_s$ } is complete to  $V''_i$  for all i = 2, ..., r
- $|V_i''| \ge c |V_i'|$ , where  $c = c(\gamma, r)$  is independent of  $\varepsilon$

To prove the theorem assuming (\*), first apply (\*) to  $V'_i = V_i$  to obtain  $v_1, v_2, ..., v_s \in V'_1$  and sets  $V'_i \subseteq V_i$ . Then apply (\*) to the sets  $V'_2, V'_3, ..., V'_r$  to obtain  $u_1, u_2, ..., u_s \in V'_2$  and sets  $V''_i \subseteq V'_i$ . Then apply (\*) to the sets  $V''_3, V''_4, ..., V''_r$  and so on.



**Lemma.** Let  $Y \subseteq V_2$  with  $|Y| \ge \varepsilon |V_2|$ . Then all but  $\varepsilon |V_1|$  vertices of  $V_1$  have  $\ge (\gamma - \varepsilon)|Y|$  neighbors in *Y*.

Proof. Let

 $X = \{v \in V_1 : v \text{ is adjacent to } < (\gamma - \varepsilon)|Y| \text{ vertices in } Y\}.$ 

Then  $|\langle X, Y \rangle| < |X|(\gamma - \varepsilon)|Y|$ . Thus  $d(X, Y) < \gamma - \varepsilon$ , and so  $|X| < \varepsilon |V_1|$ , as desired.  $\Box$ 

By the lemma, all but  $(r - 1)\varepsilon |V_1|$  vertices of  $V_1$  have  $\ge (\gamma - \varepsilon) |Y|$ neighbors in *Y* for  $Y = V'_2, V'_3, ..., V'_r$ . Pick  $v_1 \in V'_1$  with this property and let  $W_i \subseteq V'_i$  be such that  $|W_i| \ge (\gamma - \varepsilon) |V'_i|$  and  $v_1$  is complete to  $W_i$  for all i = 2, 3, ..., r. Repeat the same argument, but with  $V'_2, V'_3, ..., V'_r$  replaced by  $W_2, ..., W_r$ . We find  $v_2 \in V'_1 - \{v_1\}$  and  $Z_i \subseteq W_i$  such that  $|Z_i| \ge (\gamma - \varepsilon) |W_i|$  and  $v_2$  is complete to  $Z_i$  for all i = 2, 3, ..., r. After *s* iterations we will end up with the vertices  $v_1, v_2, ..., v_s$  and sets  $V''_2, V''_3, ..., V''_r$ . **Theorem (Triangle removal lemma)**  $\forall \varepsilon > 0 \exists \delta > 0 \exists n_0$  such that if *G* is a graph on  $\geq n_0$  vertices with at most  $\delta n^3$  triangles, then  $\exists F \subseteq E(G)$  such that  $|F| \leq \varepsilon n^2$  and  $G \setminus F$  is triangle-free.

"If a graph has  $o(n^3)$  triangles, then all triangles can be destroyed by removing  $o(n^2)$  edges"

# **Application to arithmetic progressions**

**Lemma.**  $\forall \varepsilon > 0 \exists n_0 \forall \text{graph } G \text{ on } n \ge n_0 \text{ vertices such that every edge is in exactly one triangle has <math>\le \varepsilon n^2$  edges.

**Proof.** By the Triangle removal lemma  $\exists \delta > 0 \exists n_0$  such that every graph on  $\geq n_0$  vertices with  $\leq \delta n^3$  triangles can be made  $\Delta$ free by deleting  $\varepsilon n^2$  edges. Our graph has  $|E(G)| \leq n^2$  triangles, which is  $\leq \delta n^3$  for big enough n. So if n is a big enough, then  $\exists F \subseteq E(G)$  such that  $|F| \leq \varepsilon n^2$  and  $G \setminus F$  is triangle free. Let F'consist of the edges of the unique triangle containing e, for all  $e \in F$ . Thus  $|F'| \leq 3\varepsilon n^2$  and  $E(G) \subseteq F'$ , as required.  $\Box$ 

# Definition. A corner is a triple of the form

 $\{(x, y), (x, y + d), (x + d, y)\}$  for some d, possibly negative.

**Corollary.** (Ajtai & Szemerédi 1974). Let  $A \subseteq [\{1, 2, ..., N\}]^2$ . If *A* contains no corner, then  $|A| = o(N^2)$ .

**Proof.** Consider the  $N \times N$  grid.

X = horizontal lines y = i for i = 1, 2, ..., N

Y =vertical lines x = i for i = 1, 2, ..., N

Z = slope - 1 lines y = -x + i for i = 1, 2, ..., 2N - 1

Then |X| = |Y| = N and |Z| = 2N - 1

Define a tripartite graph *G* on  $X \cup Y \cup Z$  by saying that the two lines are adjacent if their intersection belongs to *A*. Then |V(G)| = 4N - 1 and |E(G)| = 3|A|. Since *A* is corner-free, each edge belongs to a unique triangle. By the previous lemma,  $3|A| = o(N^2)$ , as desired.

**Corollary.** (Roth's theorem) If  $S \subseteq \{1, 2, ..., N\}$  contains no 3-term arithmetic progression, then |S| = o(N).

**Proof.** Let  $A = \{(x, y): x, y \in \{1, 2, ..., 2N\}, y - x \in S\}$ . If *A* has a corner, say  $(x, y), (x, y + d), (x + d, y) \in A$ , then

 $y - x \in S$ ,  $y + d - x \in S$ ,  $y - x - d \in S$ , and so *S* contains a 3term arithmetic progression. We may therefore assume that *A* has no corner. For every  $\{s, s'\} \subseteq S$  we have  $(s, s + s') \in A$ , and hence  $|A| \ge {|S| \choose 2}$ . By the previous corollary  ${|S| \choose 2} \le |A| = o(N^2)$ , and so |S| = o(N).

# **Application to property testing**

**Definition.** A graph *G* on *n* vertices is  $\varepsilon$ -far from triangle-free if for every set  $F \subseteq E(G)$  of size at most  $\varepsilon n^2$  the graph  $G \setminus F$  has a triangle.

**Remark.** We cannot hope to test whether a graph is triangle-free in constant time, but how about distinguishing triangle-free graphs from those that are  $\varepsilon$ -far from triangle-free?

**Theorem.** For every  $\varepsilon > 0$  there exists a randomized algorithm which in constant time accepts every triangle free graph and rejects every graph which is  $\varepsilon$ -far from triangle-free with probability at least 2/3.

Proof. Let  $\delta$  be as in the Triangle removal lemma. Thus if *G* is  $\varepsilon$ -far from triangle-free, then it has more than  $\delta n^3$  triangles. Pick  $\delta^{-1}$  triples of vertices uniformly independently at random. If none of those triples form a triangle, then accept the graph; otherwise reject. If the graph is a triangle-free, then it will be accepted. If it is  $\varepsilon$ -far from triangle-free, then the probability of being accepted is at most

$$\left(1 - \frac{\delta n^3}{\binom{n}{3}}\right)^{\delta^{-1}} \le \frac{1}{3}$$

if  $\delta$  is sufficiently small.